Modularity'14

Composable User-Defined Operators That Can Express User-Defined Literals

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User-Defined Operators

useful for implementing internal DSLs

- can introduce DSL-like syntax
- can be used together with other operators

An example program using internal DSLs (in Scala)

Existing User-Defined Operators

Their syntax is strictly restricted

In Scala, users can define only infix binary operators and postfix unary operators.

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Desired User-Defined Operators

Accept flexible syntax

- not only infix, prefix, postfix, or outfix
- not only unary, binary, ...
- can express literals by combining operators

```
ResultSet ids = select id from DB.students
    where entranceYear == 2013;
for (String id : ids.toList()) {
    id should match 48-13(6|7)6[0-9]{2};
    _ should __ match __ user-defined literals !
}
```

Problem: Parsing is Difficult

The grammar may be highly ambiguous

- A DSL developer cannot know all DSLs that are used together with his/her DSL
- Every operator expresses an expression
 - Expression's rule would be complex
- Especially, literal rules introduce a large number of ambiguities. (cf. regex)

Naïve Solution is Inefficient

Generate all possible parse trees and then choose the most suitable one

 Common scanner-less CFG parser takes O(n³) time if the grammar is ambiguous

* n = # of characters

- The number of trees might exponentially explode
 - * choosing the most suitable tree is difficult in a naïve way

Proposal: Using Expected Type Info for Parsing

Parser uses only operators with the expected return type

- when the parser tries to parse an expression
- an operand is parsed by operators whose return type is the operand type.
- it can reduce ambiguities since operators with the same syntax can be distinguished by types

- 1) parse a statement by the host language rules until the parser encounters an expression part
- 2) determine expected type of the next expr
- 3) pick up an operator with expected return type, and try to parse the expr by the operator's rule
- 4) if the parser encounters an operand, go to 2
- 5) if an attempt succeeds, return the result. otherwise, go to 3 and try another operator

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Parsing Precedence

For efficient parsing, we also propose to introduce parsing precedence.

- precedence rule among operators with the same return type (and the same operator precedence)
- which operator is chosen if an expr is ambiguous
- can remove all ambiguities, but may change the grammar

Efficiency

O(n) time for practical grammar

- ambiguities are removed by
 - * using static types as non-terminal symbols
 - * parsing precedence
- using memoization
 - * for reducing the cost of backtrack
 - * packrat parsing supporting left-recursion

Benefits

Operators can express literals

- literal is an expression with special whitespace rule
- literals overloads token rules

Drawbacks

Limited places where operators are available

- only in expressions whose expected type is statically determined before parsing
- depends on the host language
- e.g. the receiver of a method call in Java

not available	available
<u>a</u> = b	a = <u>b</u>
<u>a</u> .method(x, y)	a.method(<u>x</u> , <u>y</u>)
<u>a</u> .field	<u>e;</u>
(Type) <u>a</u>	return <u>a</u> ;
	•••

Implementation: ProteaJ

A subset of Java with

- flexible user-defined operators
- provides a module system for operators
- does not support generics

Source and test programs are available from: https://github.com/csg-tokyo/proteaj.git

Operators in ProteaJ

Operator syntax = { name | operand }+

- Not only infix, prefix, unary, binary, ...
- Expressiveness is equivalent to PEG

Extra features

- operator precedence and associativity
- two whitespace rules : expression / literal level

Experiment

Experiment Environment CPU: 2.67 GHz Core i5 Memory: 8GB OS: OpenSUSE 12.1 Java: OpenJDK 1.7.0

Our compiler vs. JSGLR parser

JSGLR: well-known scanner-less CFG parser It can generate all possible trees

Problem settings

- grammar: arithmetic operators + file path
- input: a/a/a/.../a (input size = # of a)

note: In the ProteaJ experiment, the input is embedded in a minimal program. In the case of JSGLR, it parsed as is.

Result (semi-log graph)



Compilation time by ProteaJ (linear-scale)



Related Work: User-Defined Operators

Mixfix operators

- a class of user-defined operators
- only infix, prefix, postfix, or outfix
- Coq, Agda, Pure, OBJ3, Isabelle, ...

Mixfix operators + implicit (empty) operators

- mixfix + operator having no name
- poorly supports user-defined literals
- OBJ3, Isabelle

Related Work: Parsing

CFG parser + type-based disambiguation

- generate all ASTs => type check
- inefficient for highly ambiguous grammar
- Metaborg, Agda, OBJ3, Isabelle, ...

Type-oriented island parsing [Silkenson '12]

- bottom-up parsing using type information
- cannot define new (complex) literals

Conclusion

Parsing method for flexible operators

- using expected type information
- precedence rule: parsing precedence
- O(n) parse time for practical grammar

Benefits

• Operators can express literals

Drawbacks

• Limited places where operators are available

Efficiency

O(n) time for practical grammar

- ambiguities are removed by
 - * using static types as non-terminal symbols
 - * parsing precedence
- operators ≒ PEG including left-recursion
 - * operator name \Rightarrow terminal
 - * return type, operand type ≒ non-terminal
 - * parsing precedence ≒ ordered choice

Parsing Precedence

For efficient parsing, we also propose to introduce parsing precedence.

- precedence rule among operators with the same return type (and the same operator precedence)
- which operator is chosen if an expr is ambiguous
- parsing precedence ≒ ordered choice rule
- it is declared by programmers

User-Defined Literals

Protean operators can express literals



Motivation: Internal DSL

DSL implemented as a **library**

- It can be used as a part of its host language
- e.g. parser combinator, OR mapper

DSL (Domain Specific Language)

- specialized language for a specific purpose
- e.g. yacc, SQL

Pros/Cons of Internal DSL

Pros: Composability

It can be used together with other DSLs

Cons: Syntax

The syntax is restricted by its host language

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Goal: Composable Syntax Extension

Enabling a DSL to introduce its own syntax

- the syntax is not restricted by the host lang.
- the syntax includes literal-level syntax

Without breaking composability

- multi-DSLs can be used together safely
- without critical penalty of compilation time

Proposal: Protean Operators

- A class of user-defined operators
 - consist of names and operands
 - * not only infix, prefix, postfix, and outfix
 - overloaded by its return type
 - * an operator is available only at an expression where the return type is expected
 - have a special rule: parsing precedence
 - * Programmers should declare the precedence among operators with the same return type

Protean Operators Introduce DSL Syntax

DSL syntax can be expressed by protean operators !

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ResultSet ids = select id from DB.students
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for (String id : ids.toList() ) {
    id should match 48-13(6|7)6[0-9]{2};
    _ should __ match __ regular expression literals
```

Protean Operators Introduce DSL Syntax

DSL syntax can be expressed by protean operators !



Protean Operators are Composable

Compiler can distinguish operators by types even if they have the same syntax !

```
ResultSet ids = select id from DB.students
    where entranceYear == 2013;
for (String id : ids.toList() ) {
    id should match 48-13(6|7)6[0-9]{2};
}
2013 is an integer literal
    because int is expected.
This part is parsed by regex operators
    because Regex is expected here !
```

We developed a parsing method that uses expected type information.

- 1) parse a statement by the host language rules until the parser encounters an expression part
- 2) determine the expected type of the next expression
- 3) parse the expression by the operators that return the expected type
- 4) if the parser encounters an operand, go to 2

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Implementation: ProteaJ

A subset of Java + protean operators

- provides module system for operators
- not supports generics

Source and test programs are available from: <u>www.csg.ci.i.u-tokyo.ac.jp/~ichikawa/ProteaJ.tar.gz</u>

Expressiveness of Protean Operators

Pros: They can express any PEGs

- non-terminal => static type
- PEG (Parsing Expression Grammar) is a type of formal grammar like CFG

Cons: They cannot express declarations

• They do not use meta-programming

Efficiency of Our Parsing Method

O(n) for practical grammar

- n: input source length (# of letters)
- use memoization to reduce back-track cost

Naive method is inefficient

- generate all possible ASTs and then choose most suitable one by using types
- parser that can generate all possible ASTs is inefficient against highly ambiguous grammar

Experiment

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Problem settings

- grammar: arithmetic operators + file path
- input: a/a/a/.../a (input size = # of a)

note: The input for our compiler is more complex since it must be a valid ProteaJ program.

Result (semi-log graph)



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Conclusion

Protean operators

- expressiveness is equivalent to PEG
- multiple operators can be used safely

Parsing method

- uses expected type information
- O(n) for practical grammar